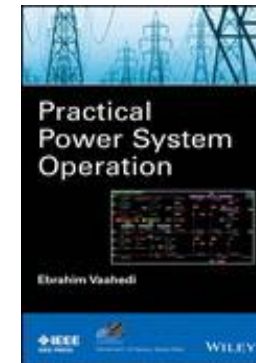
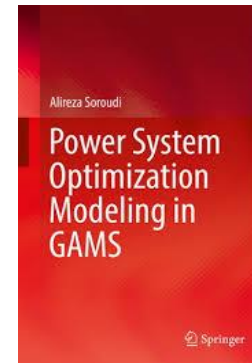
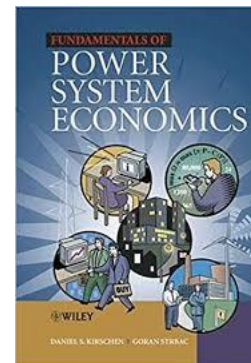
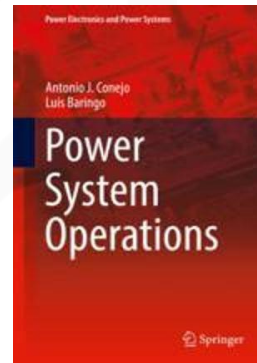
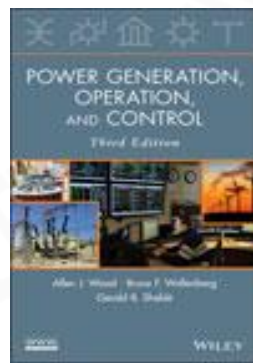
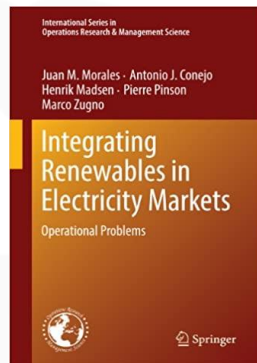


Acknowledgement

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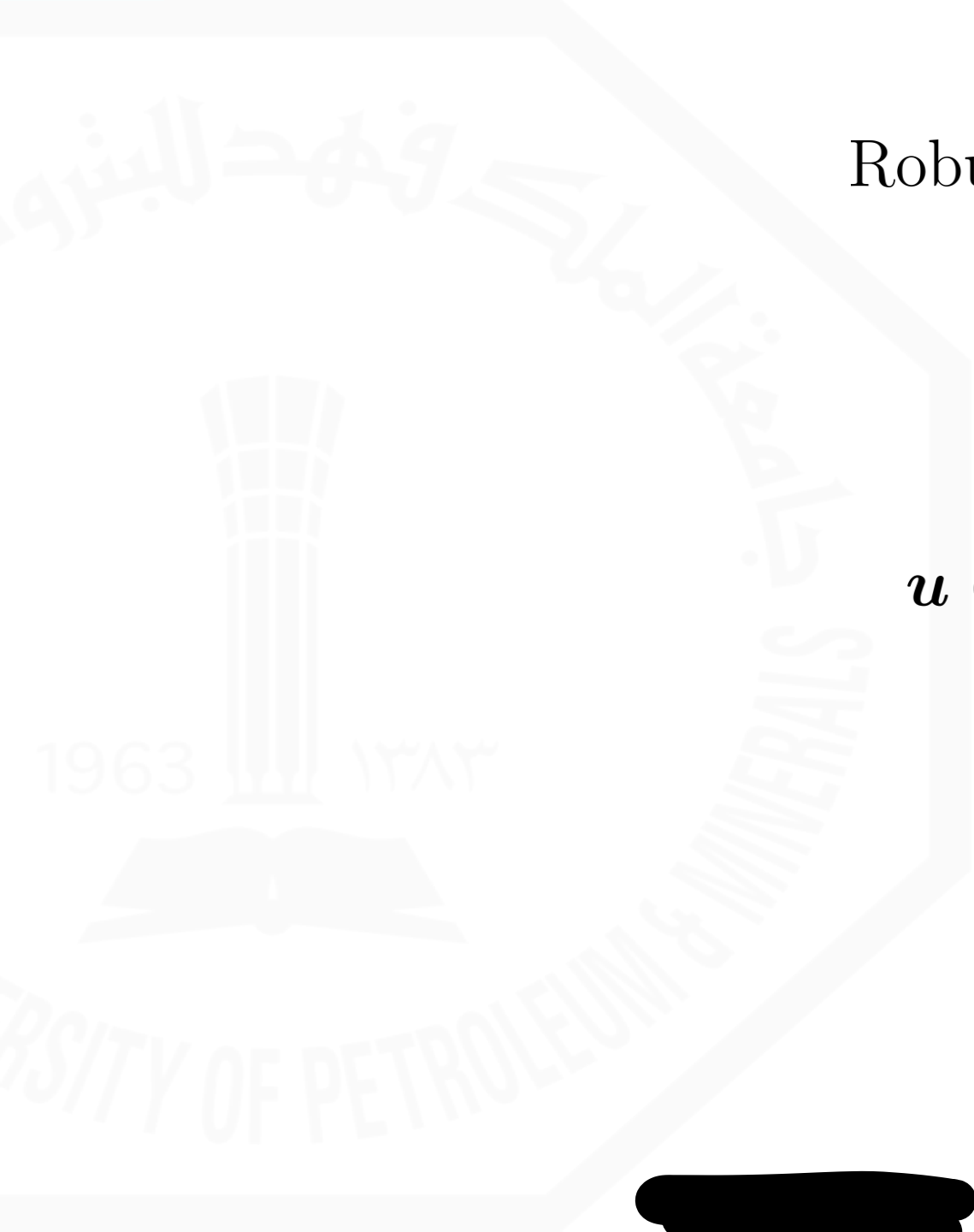
What

Column-and-Constraint Generation (CCG) algorithm



Robust set

$$u \in \mathcal{U}$$



ARO

$$\begin{aligned} \min_x \\ \text{s.t.} \\ h(x) = 0 \\ g(x) \leq 0 \end{aligned}$$

$$\begin{aligned} \max_u \\ \text{s.t.} \\ u \in \mathcal{U} \end{aligned}$$

$$\min_y \quad c_I^\top x + [c_O(x, u)]^\top y$$

$$\begin{aligned} \text{s.t.} \quad A(x, u) \cdot y = b(x, u) : \quad \lambda \\ D(x, u) \cdot y \geq e(x, u) : \quad \mu \end{aligned}$$

$c_I^\top x$ is a constant

Third-level problem

$$\begin{aligned} \min_y & \quad [c_0(x, u)]^\top y \\ \text{s.t.} & \quad A(x, u) \cdot y = b(x, u) : \lambda \\ & \quad D(x, u) \cdot y \geq e(x, u) : \mu \end{aligned}$$

Dual of the third-level problem

$$\begin{aligned} \min_y \quad & [c_0(x, u)]^\top y \\ \text{s.t.} \quad & A(x, u) \cdot y = b(x, u) : \lambda \\ & D(x, u) \cdot y \geq e(x, u) : \mu \end{aligned}$$

$$\begin{aligned} \max_{\lambda, \mu} \quad & [b(x, u)]^\top \lambda + [e(x, u)]^\top \mu \\ \text{s.t.} \quad & [A(x, u)]^\top \lambda + [D(x, u)]^\top \mu = c_0(x, u) \\ & \lambda : \text{ free} \\ & \mu \geq 0 \end{aligned}$$

Second-level problem and dual of the third-level problem merged

$$\begin{aligned} \max_{u, \lambda, \mu} & \quad [b(x, u)]^\top \lambda + [e(x, u)]^\top \mu \\ \text{s.t.} & \quad u \in \mathcal{U} \\ & \quad [A(x, u)]^\top \lambda + [D(x, u)]^\top \mu = c_0(x, u) \\ & \quad \lambda : \text{ free} \\ & \quad \mu \geq 0 \end{aligned}$$

Subproblem: $x = x^{(\nu-1)}$ fixed

Subproblem: $\mathbf{x} = \mathbf{x}^{(\nu-1)}$ fixed

$$\max_{\mathbf{u}, \boldsymbol{\lambda}, \boldsymbol{\mu}} \quad [\mathbf{b}(\mathbf{x}^{(\nu-1)}, \mathbf{u})]^\top \boldsymbol{\lambda} + [\mathbf{e}(\mathbf{x}^{(\nu-1)}, \mathbf{u})]^\top \boldsymbol{\mu}$$

$$\text{s.t.} \quad \mathbf{u} \in \mathcal{U}$$

$$[\mathbf{A}(\mathbf{x}^{(\nu-1)}, \mathbf{u})]^\top \boldsymbol{\lambda} + [\mathbf{D}(\mathbf{x}^{(\nu-1)}, \mathbf{u})]^\top \boldsymbol{\mu} = \mathbf{c}_0(\mathbf{x}^{(\nu-1)}, \mathbf{u})$$

$$\boldsymbol{\lambda} : \text{ free}$$

$$\boldsymbol{\mu} \geq \mathbf{0}$$

$$\Downarrow$$
$$\mathbf{u}^{(\nu)}, \boldsymbol{\lambda}^{(\nu)}, \boldsymbol{\mu}^{(\nu)}$$

Upper bound: \mathbf{x} arbitrarily fixed
(\mathbf{x} to be fixed to minimize the objective function)

$$\text{UB}_\nu = \mathbf{c}_1^\top \mathbf{x}^{(\nu-1)} + [\mathbf{b}(\mathbf{x}^{(\nu-1)}, \mathbf{u}^{(\nu)})]^\top \boldsymbol{\lambda}^{(\nu)} + [\mathbf{e}(\mathbf{x}^{(\nu-1)}, \mathbf{u}^{(\nu)})]^\top \boldsymbol{\mu}^{(\nu)}$$

Master problem: $u = u^{(k)}, k = 1, \dots, \nu$, fixed

Master problem: $u = u^{(k)}, k = 1, \dots, \nu$, fixed

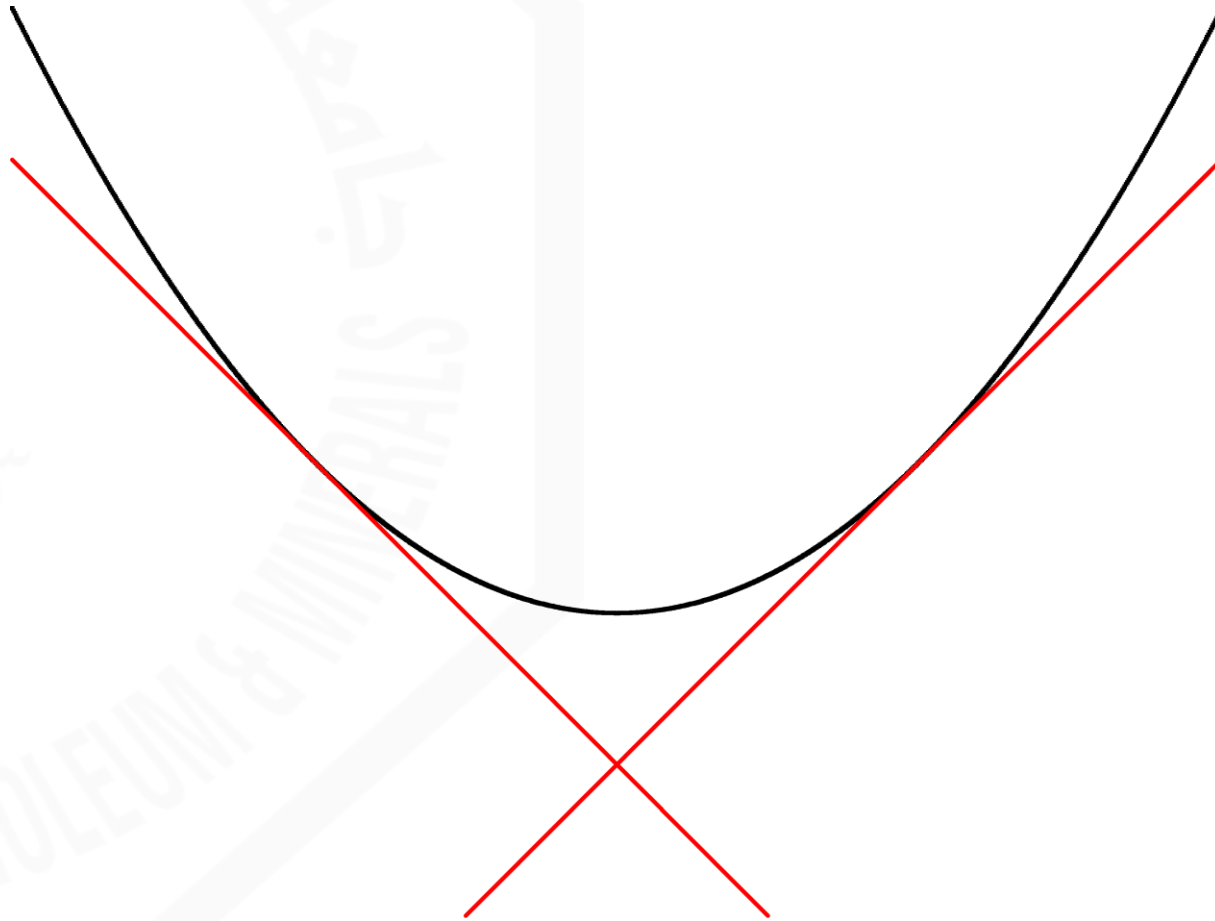
$$\begin{array}{ll} \min_{x, \eta, y^{(k)}, k=1, \dots, \nu} & c_I^\top x + \eta \\ \text{s.t.} & h(x) = 0 \\ & g(x) \leq 0 \\ & \eta \geq [c_0(x, u^{(k)})]^\top y^{(k)} & k = 1, \dots, \nu \\ & A(x, u^{(k)}) \cdot y^{(k)} = b(x, u^{(k)}) & k = 1, \dots, \nu \\ & D(x, u^{(k)}) \cdot y^{(k)} \geq e(x, u^{(k)}) & k = 1, \dots, \nu \end{array}$$

↓

$$x^{(\nu)}, \eta^{(\nu)}$$

($\& y^{(k)}, k = 1, \dots, \nu$)

Lower bound: bounded from below ($\eta \geq [\mathbf{c}_0(\mathbf{x}, \mathbf{u}^{(k)})]^\top \mathbf{y}^{(k)}$)



Lower bound: bounded from below ($\eta \geq [\mathbf{c}_0(\mathbf{x}, \mathbf{u}^{(k)})]^\top \mathbf{y}^{(k)}$)

$$\text{LB}_\nu = \mathbf{c}_I^\top \mathbf{x}^{(\nu)} + \eta^{(\nu)}$$

1. Solve the SP and get $\mathbf{u}^{(\nu)}$ (with previous or initial $\mathbf{x}^{(\nu-1)}$)
2. Compute the upper bound UB_{ν}
3. Solve the MP and get $\mathbf{x}^{(\nu)}$
4. Compute the lower bound LB_{ν}
5. If $UB_{\nu} - LB_{\nu}$ is small enough, stop; otherwise continue in 1

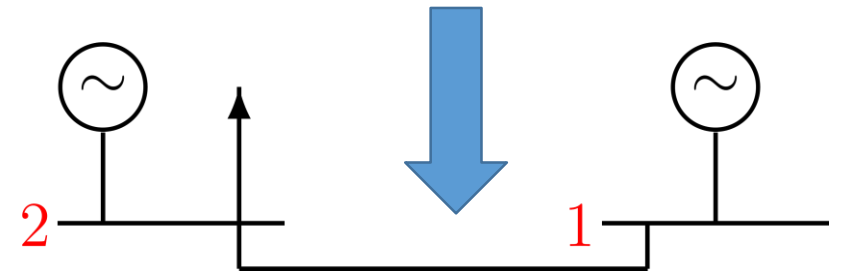


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Simple adaptive robust optimization example

By utilizing GCC approach, implement the following Adaptive Robust Optimization using GAMS



A transmission expansion planner has the option of building just one of two alternatives transmission lines of capacity either 2 or 4 to connect electrical nodes 1 and 2.

The per unit building cost for the transmission lines is 1.

ARO

Cost = 10

Cost = 5



At node 1, a generator of unlimited capacity and per unit production cost of 5 is available, while in node 2 another generator of unlimited capacity and per unit production cost of 10 is available.

The demand is located at node 2 and can take solely two values, either 3 or 5.

ARO



$$\min_{x \in \{2,4\}} \quad \max_{d \in \{3,5\}} \quad \min_{y_1, y_2 \geq 0}$$

s.t.

$$z = 1x + 5y_1 + 10y_2$$

$$y_1 + y_2 = d$$

$$y_1 \leq x$$

where x represents the capacity of the transmission line to be built,
 d the unknown demand, and

y_1 and y_2 , the production of the units at nodes 1 and 2, respectively.

ARO

$$\min_{\mathbf{x} \in \mathcal{X}} \max_{\mathbf{u} \in \mathcal{U}} \min_{\mathbf{y} \in \mathcal{Y}(\mathbf{x}, \mathbf{u})} f(\mathbf{x}, \mathbf{y}, \mathbf{u})$$

On the computational side,

we note that **tri-level problems require a specific solution algorithm** that generally relies on

merging the second- and third-level problems using duality theory and

decomposing the resulting bi-level problem through a **column-and-constraint generation** algorithm.

Thank you!

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